

FIG 1

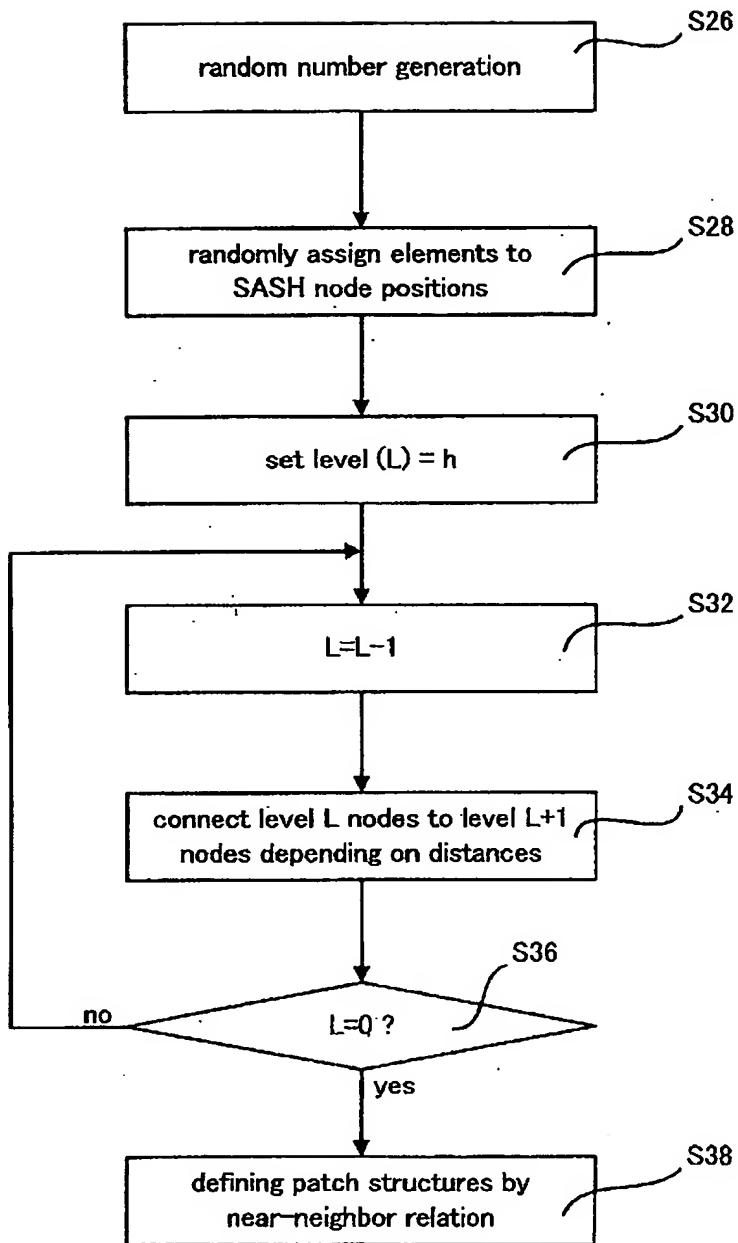


FIG 2

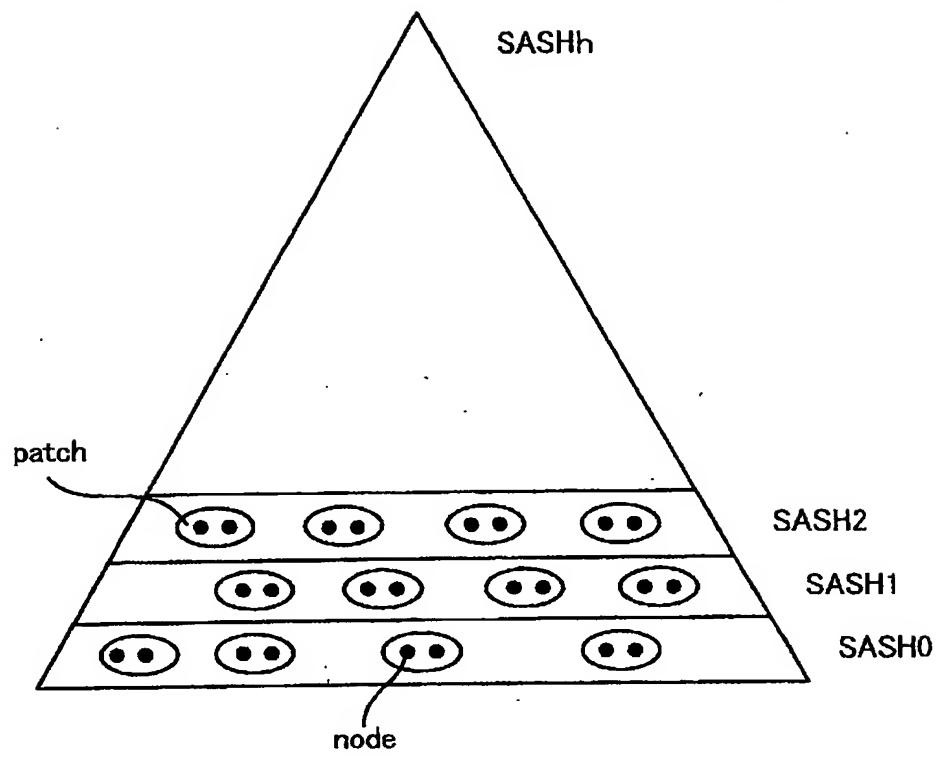


FIG 3

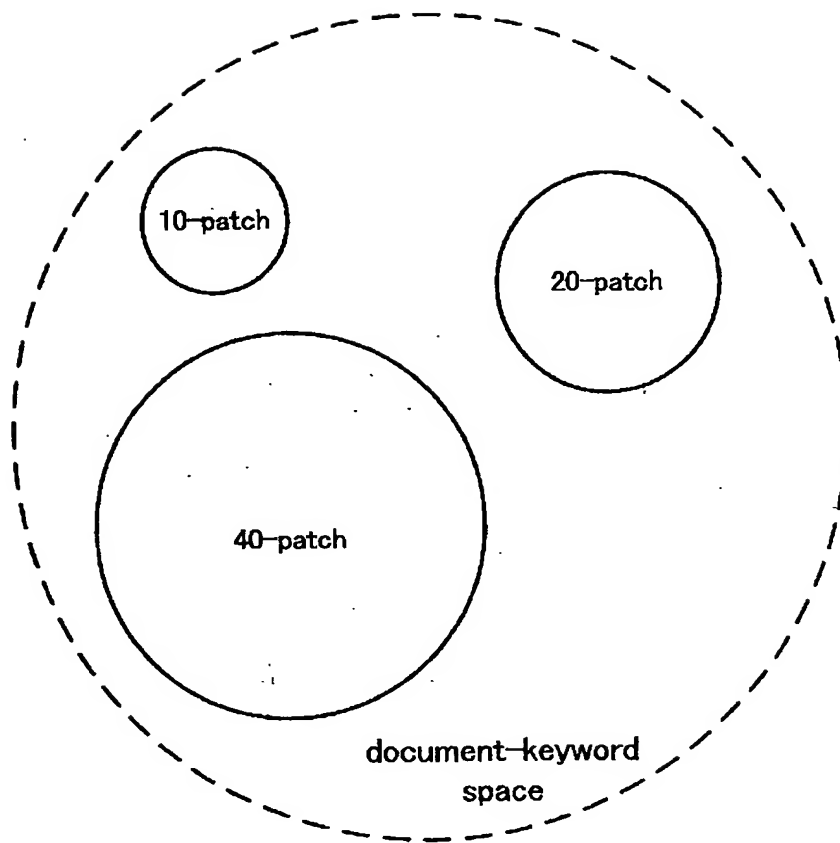
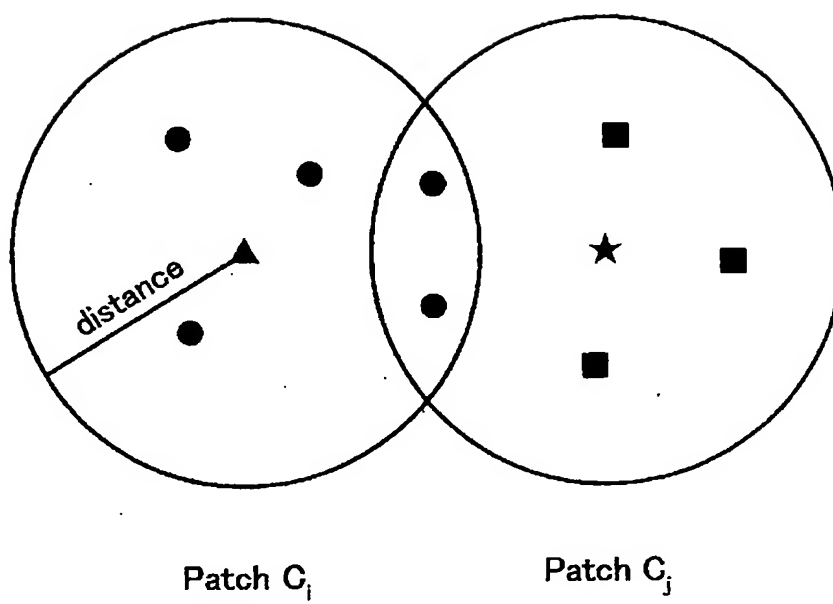


FIG 4



$$\text{CONF}(C_i, C_j) = 2/5 = 40\%$$

FIG 5

Profile (query  $q$ ; maximum patch size  $m$ ): SCONE list SCONFL

{Let  $QNL$  be the  $m$ -patch precomputed for query  $q$ .}

{Let  $NNL$  be a list of the  $m$ -patches precomputed for every element of  $QNL$ .}

{Initially,  $w.count = \emptyset$  is assumed for every element  $v$  in the data set.}

1.  $score \leftarrow 0$ ;  
 {Initially, no query neighbors are in the current patch.}  
 for  $i = 1$  to  $m$  do
2.  $QNL[i].count \leftarrow 0$ ;  
 end for  
 for  $i = 1$  to  $m$  do  
 {Retrieve the number of times  $QNL[i]$  has been encountered as an external neighbor so far.}
3.  $score \leftarrow score + QNL[i].count$ ;  
 {Indicate that henceforth  $QNL[i]$  is in the current  $i$ -patch.}
4.  $QNL[i].count \leftarrow present$ ;  
 for  $j = 1$  to  $i - 1$  do
5.  $w \leftarrow NNL[j, i]$ ;  
 if  $w.count = present$  then
6.  $score \leftarrow score + 1$ ;
7.  $w.count \leftarrow w.count + 1$ ;  
 end if
8.  $w \leftarrow NNL[i, j]$ ;  
 if  $w.count = present$  then
9.  $score \leftarrow score + 1$ ;
10.  $w.count \leftarrow w.count + 1$ ;  
 end if
11.  $w \leftarrow NNL[i, i]$ ;  
 if  $w.count = present$  then
12.  $score \leftarrow score + 1$ ;
13.  $w.count \leftarrow w.count + 1$ ;  
 end if
14.  $SCONFL[i] = score/i^2$ ;  
 end for  
 {Reset the counts to their default value.}  
 for  $i = 1$  to  $m$  do
15.  $QNL[i].count \leftarrow \emptyset$ ;  
 end for

FIG 6

SASH #	patch list	CONFL	SCONFL	RSCONFL
0	NN( $R_{1,0},m$ ), ...	...	...	...
1	NN( $R_{1,0},m$ ), ...	...	...	...
...	...	...	...	...
...	...	...	...	...
...	...	...	...	...
h	NN( $R_{1,0},m$ ), ...	...	...	...

**FIG 7**

```

RefineProfile (query  $q$ ;
    inner patch size  $k_I$ ;
    outer patch size  $k_O$ ): reordered query  $k_I$ -patch  $RQNL$ 
{Let  $QNL$  be the  $k_O$ -patch precomputed for query  $q$ .}
{Let  $NNL$  be a list of the  $k_O$ -patches precomputed for every element of  $QNL$ .}
{Initially,  $v.inpatch = false$  is assumed for every element  $v$  in the data set.}
{Identify the inner patch members.}
for  $i = 1$  to  $k_I$  do
1.    $QNL[i].inpatch \leftarrow true$ ;
end for
{Initialize the confidence value  $CONF_c$  of every patch element to zero.}
for  $i = 1$  to  $k_O$  do
2.    $CONF_c[i] \leftarrow 0$ ;
end for
{For each element of the outer patch, count the number of elements
of their  $k$ -nearest-neighbor sets shared with that of  $q$ .}
for  $i = 1$  to  $k_O$  do
    for  $j = 1$  to  $k_I$  do
3.        $w \leftarrow NNL[i, j]$ ;
        if  $w.inpatch = true$  then
4.            $CONF_c[i] \leftarrow CONF_c[i] + 1$ ;
        end if
    end for
5.    $CONF_c[i] \leftarrow CONF_c[i] / k_O$ ;
end for
{Reorder the outer patch elements according to their confidence values, from highest to lowest.}
6.  $RQNL \leftarrow sort(QNL, CONF_c, k_O)$ ;
{Reset the patch membership indicators to their default values.}
for  $i = 1$  to  $k_I$  do
7.    $QNL[i].inpatch \leftarrow false$ ;
end for

```

FIG 8



*PatchCluster* (data set  $S$ ;

RSCM parameters  $a, b, m = \varphi(b)$ ;

Thresholds  $\alpha, \beta, \gamma, \delta$ ): query cluster graph  $G$

1. Randomly partition the set  $S$  into subsets  $S_t$  of approximate size  $\frac{|S|}{3^t}$ , for  $0 \leq t \leq h = \lceil \log_3 |S| \rceil$ .

2. For all  $0 \leq t \leq h$  do:

(a) For every element  $v \in S_t$ , compute nearest-neighbor patches  $NN(R_t, v, m)$ , where  $R_t = \bigcup_{i \geq t} S_i$ .

(b) For each element  $v_{t,i} \in S_t$ , compute the optimal query cluster size  $k(v_{t,i})$  maximizing  $RSCONF(NN(R_t, v_{t,i}, k), \varphi)$ , for values of  $k$  between  $a$  and  $b$

The ranked collection of patches

$$C_t = \langle C_{t,i} | i < j \Rightarrow RSCONF(C_{t,i}, \varphi) \geq RSCONF(C_{t,j}, \varphi) \rangle$$

form the candidates for the query clusters associated with sample  $R_t \subseteq S$ , where  $C_{t,i} = NN(R_t, v_{t,i}, k(v_{t,i}))$  and  $C_{t,j} = NN(R_t, v_{t,j}, k(v_{t,j}))$ .

(c) Let  $Q_t$  be a list of patches of  $C_t$  that have been confirmed as query clusters of  $R_t$ . Initially,  $Q_t$  is empty.

(d) For all  $1 \leq i \leq |C_t|$  do:

i. If  $RSCONF(C_{t,i}, \varphi) < \alpha$ , then break from the loop.

ii. For all  $w \in C_{t,i}$  do: if  $NN(R_t, w, k) \notin |C_t|$  for any value of  $k$ , or failing that, if  $\max\{CONF(NN(R_t, w, k), C_{t,i}), CONF(C_{t,i}, NN(R_t, w, k))\} < \beta$ , then add  $C_{t,i}$  to the list  $Q_t$ .

3. Let  $h'$  be the largest index for which  $|Q_{h'}| > 0$ . Let  $\{C_{t,j}\}$  be the set of patches comprising  $Q_t$ , where  $C_{t,j} = NN(R_t, q_{t,j}, k(q_{t,j}))$ , for all  $0 \leq t \leq h'$ . Initialize the node set of the query cluster graph  $G$  to contain these patches, one patch per node.

4. For all  $\delta \leq t \leq h'$ , all  $1 \leq j \leq |Q_t|$ , and all  $\max\{0, t - \delta\} \leq s \leq t$ , do:

(a) Compute  $C'_{t,j} = NN(R_s, q_{t,j}, 2^{t-s}k(q_{t,j}))$ .

(b) For all  $1 \leq i \leq |Q_s|$ , if  $C_{s,i} \neq C'_{t,j}$  and  $\max\{CONF(C_{s,i}, C'_{t,j}), CONF(C'_{t,j}, C_{s,i})\} \geq \gamma$ , then introduce the edges  $(C_{s,i}, C_{t,j})$  and  $(C_{t,j}, C_{s,i})$  into  $G$ , with weights  $CONF(C_{s,i}, C'_{t,j})$  and  $CONF(C'_{t,j}, C_{s,i})$ , respectively.

FIG 9

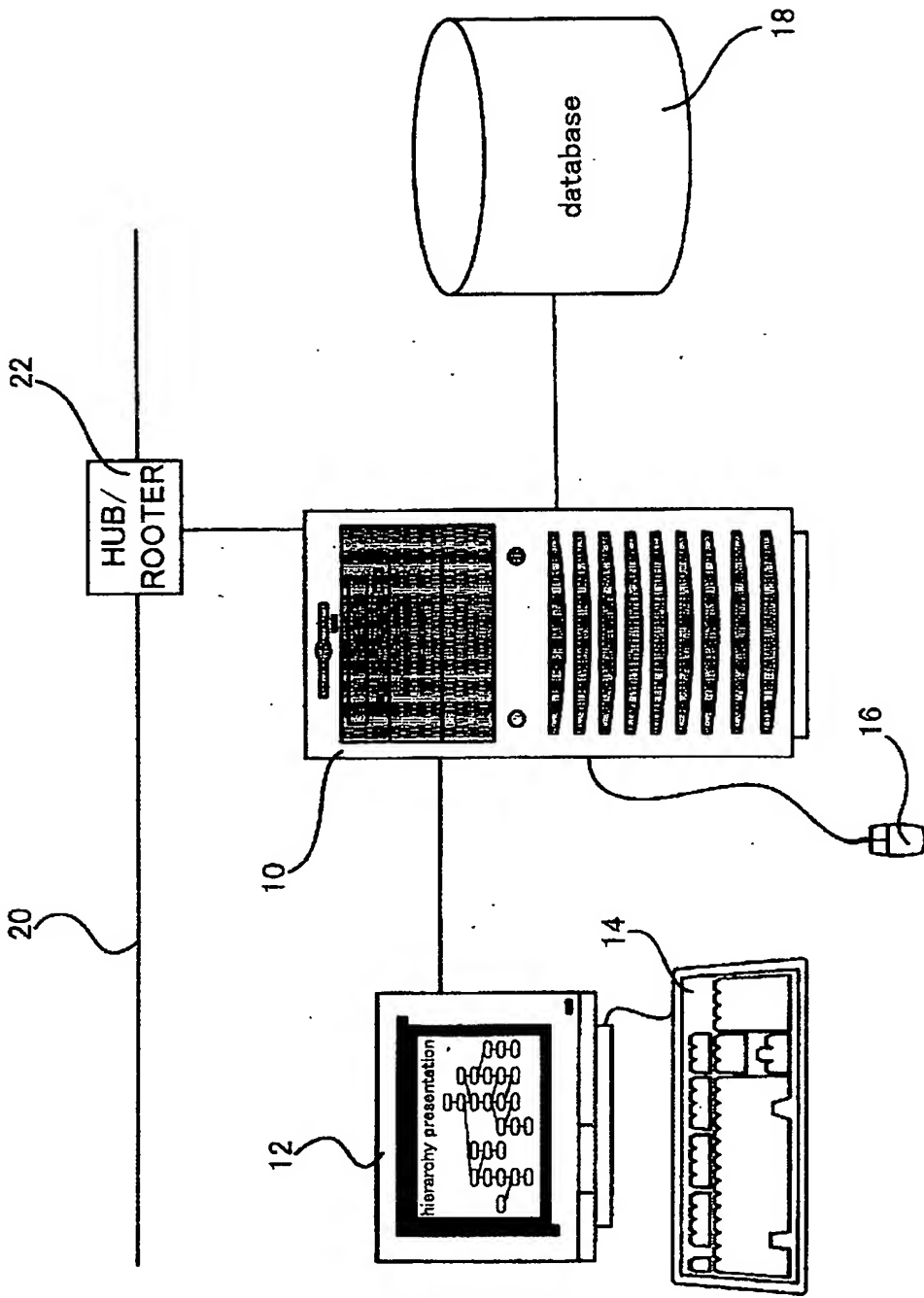


FIG 10

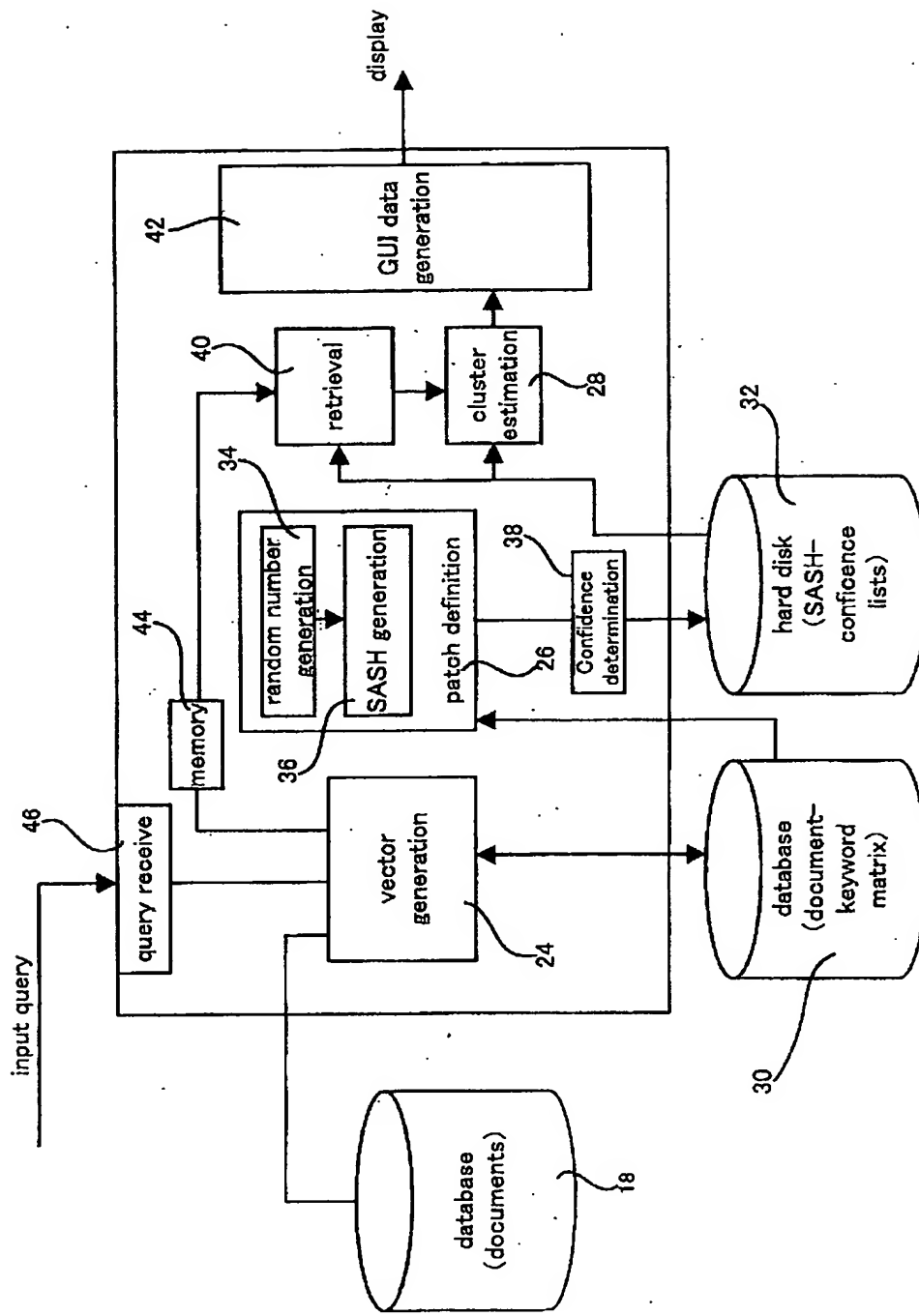


FIG 11

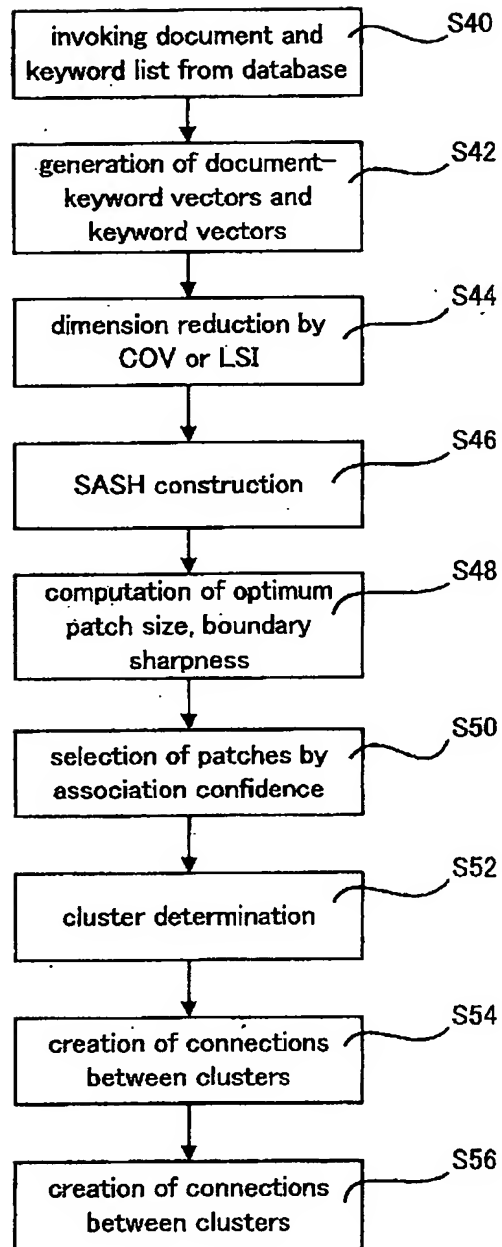


FIG 12

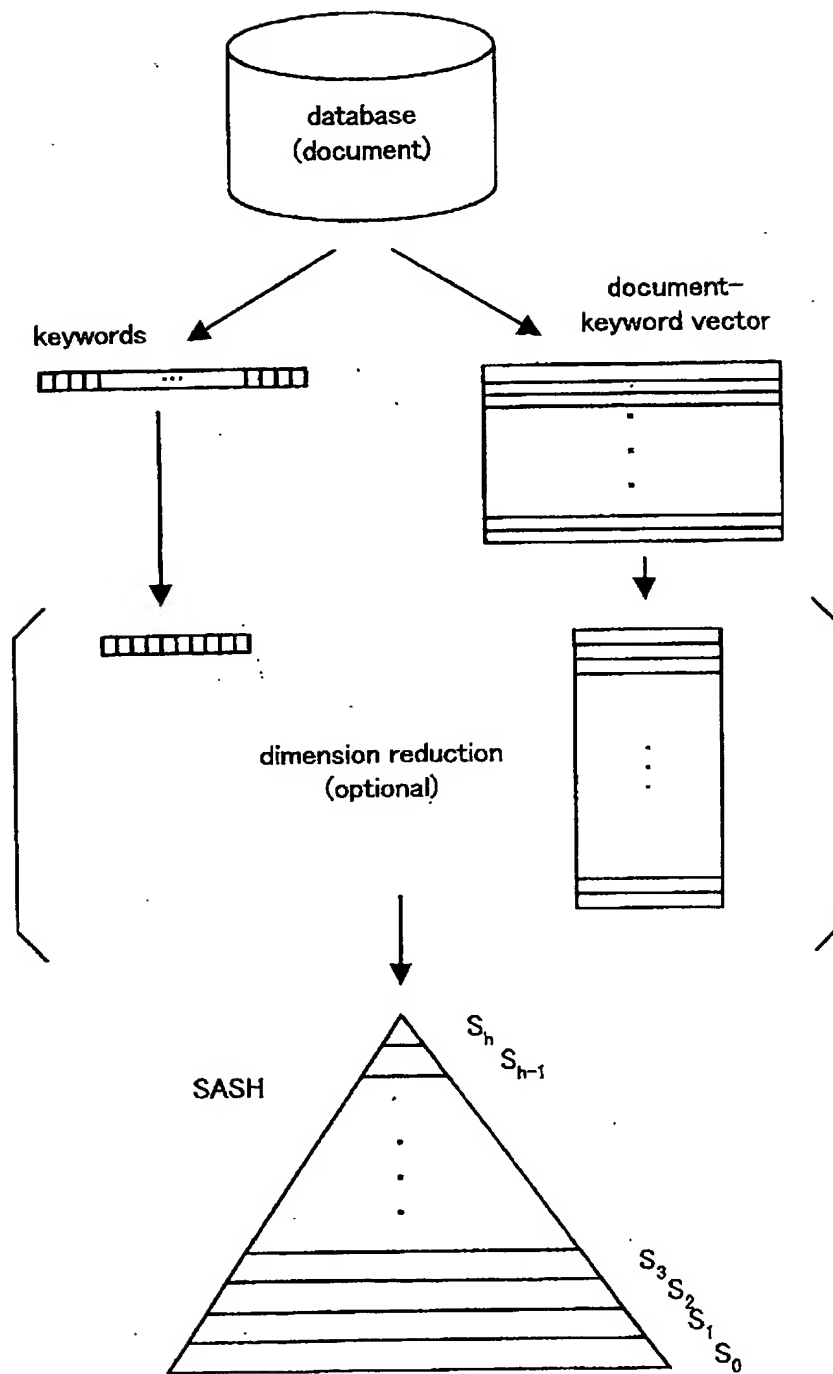
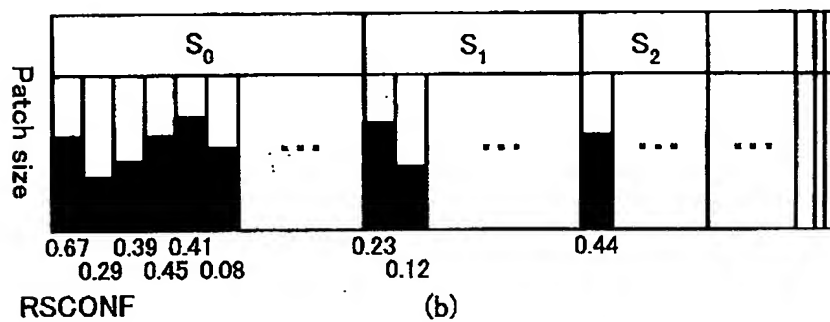


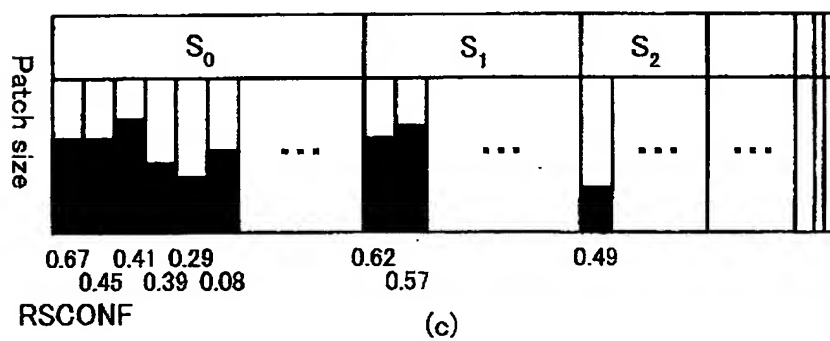
FIG13

$S_0$						$S_1$		$S_2$			
$NN(R_0, V_{0,0}, m)$	$NN(R_0, V_{0,1}, m)$	$NN(R_0, V_{0,2}, m)$	$NN(R_0, V_{0,3}, m)$	$NN(R_0, V_{0,4}, m)$	$NN(R_0, V_{0,5}, m)$	$NN(R_1, V_{0,0}, m)$	$NN(R_1, V_{0,1}, m)$	$NN(R_2, V_{0,0}, m)$	...	...	
...						...		...	...	...	

(a)



(b)



(c)

FIG 14

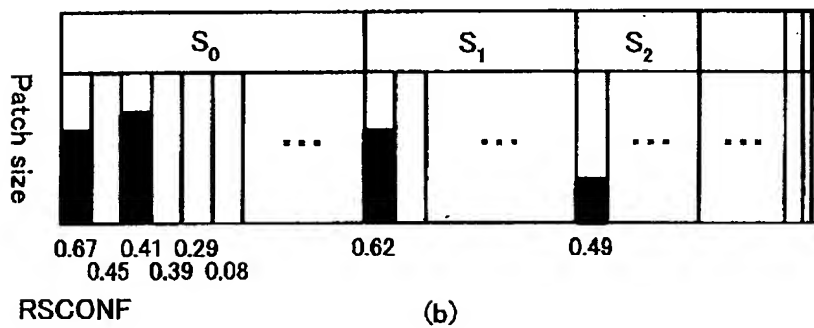
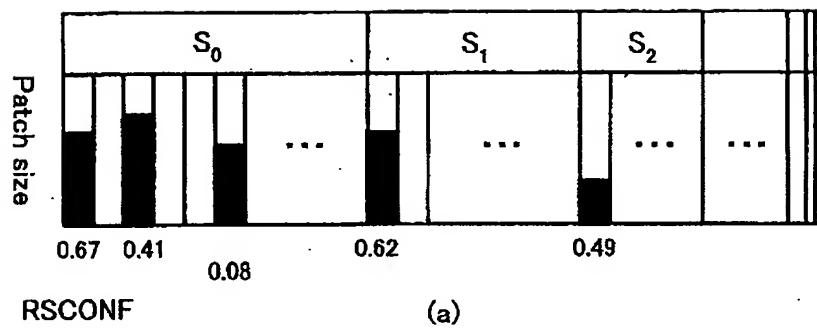


FIG 15

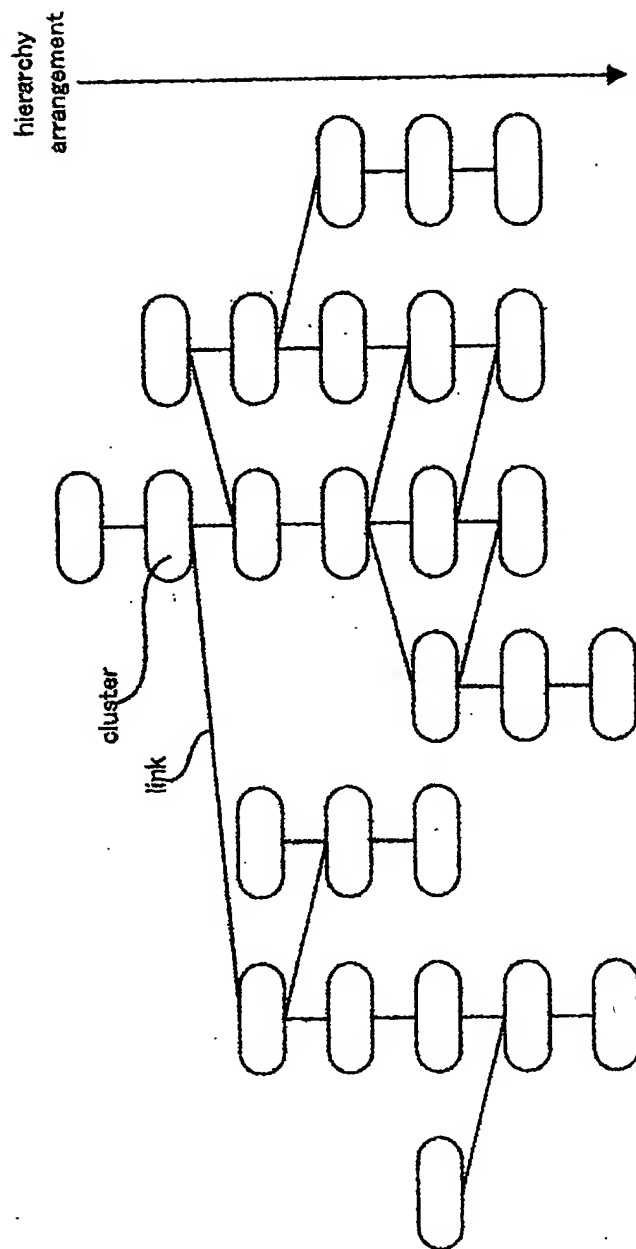


FIG 16



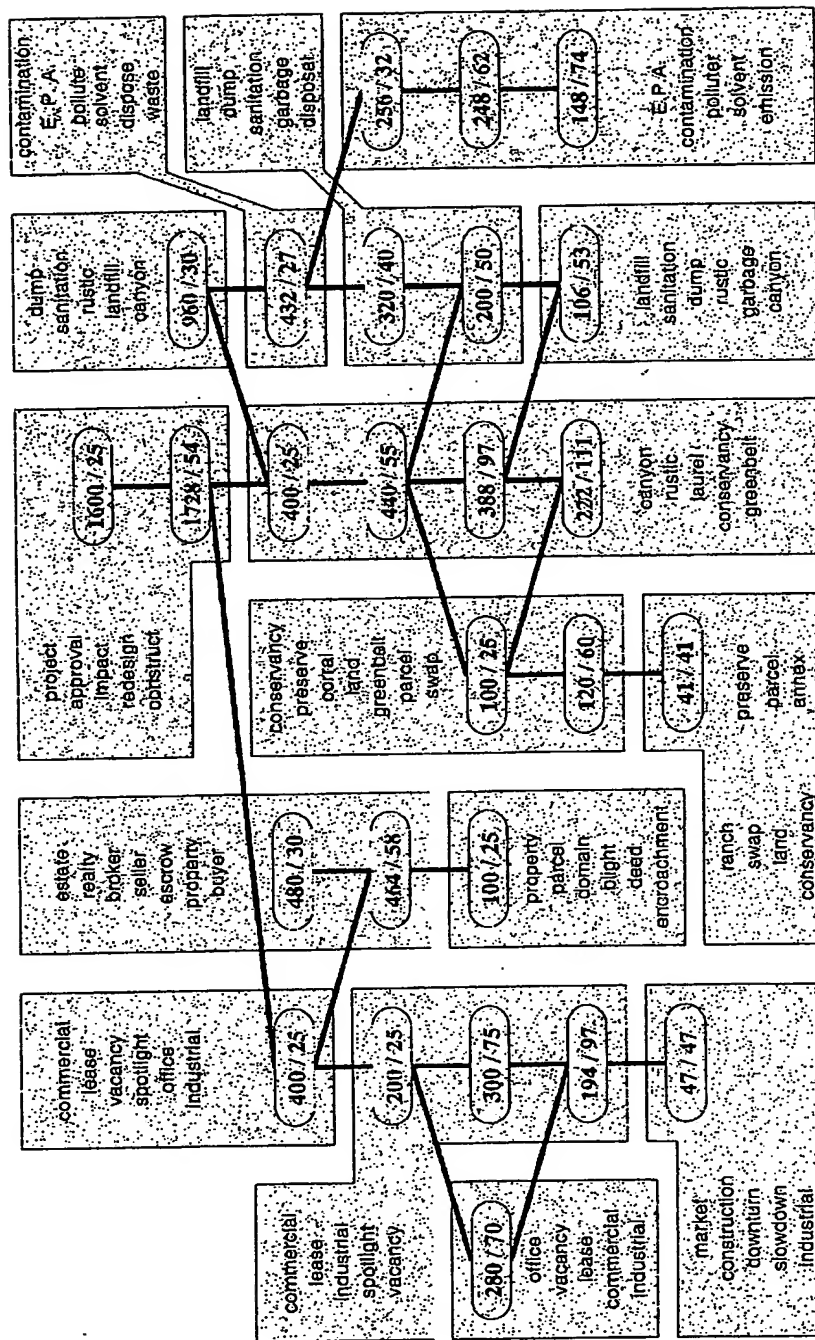


FIG 17

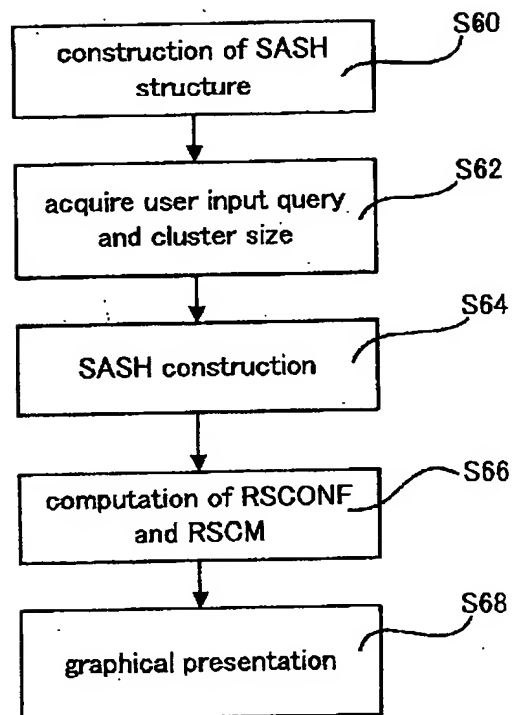


FIG 18

Patches including queried nodes

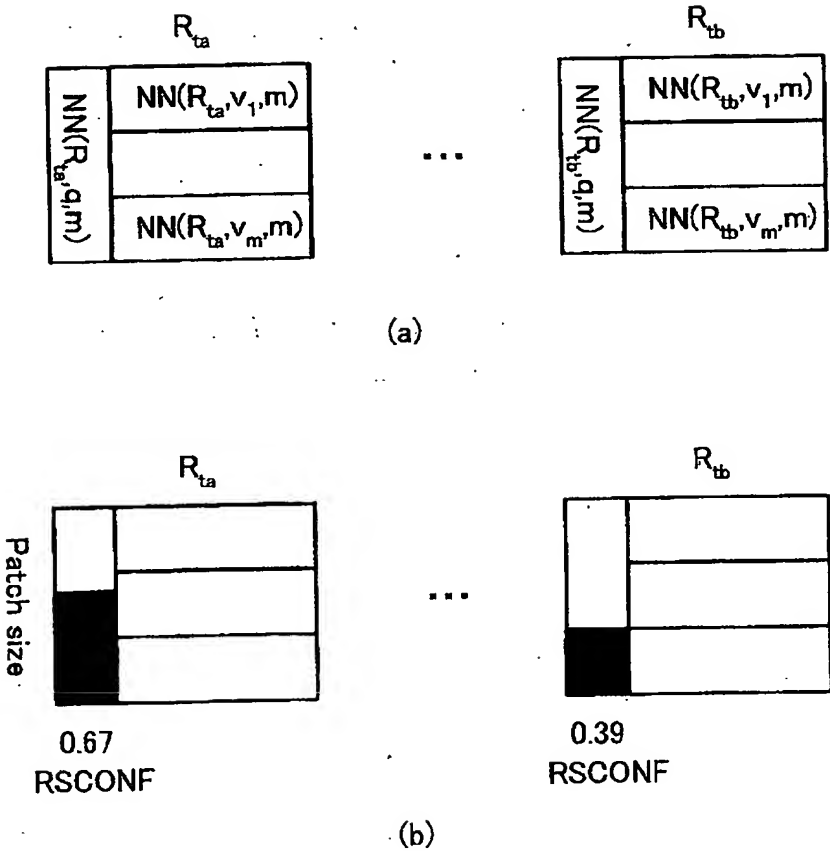


FIG 19

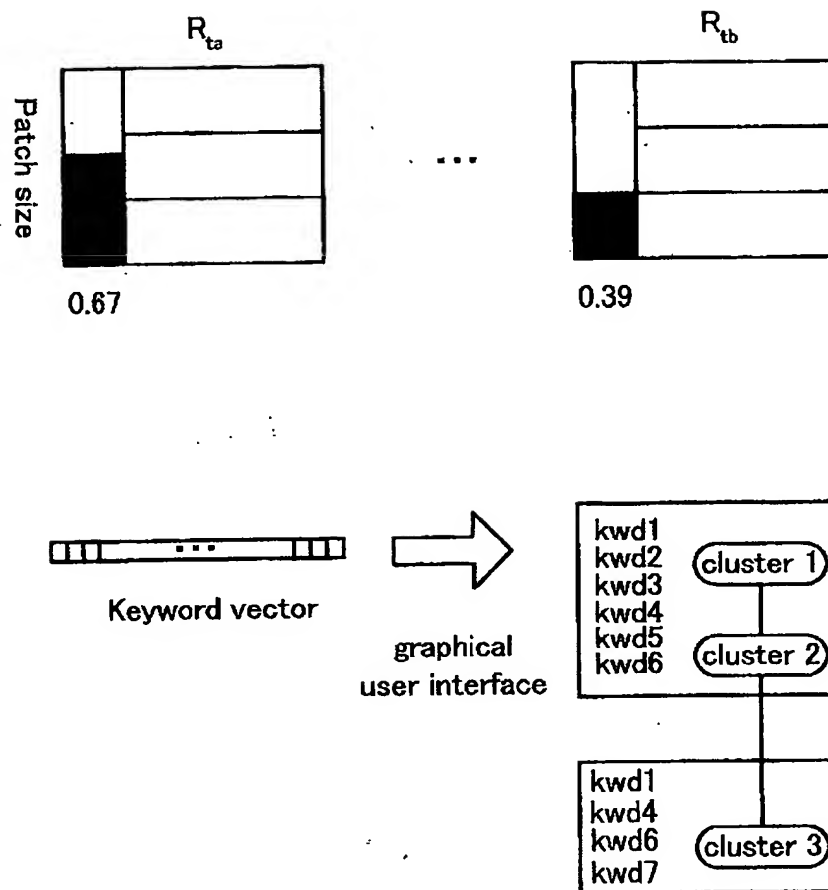


FIG 20

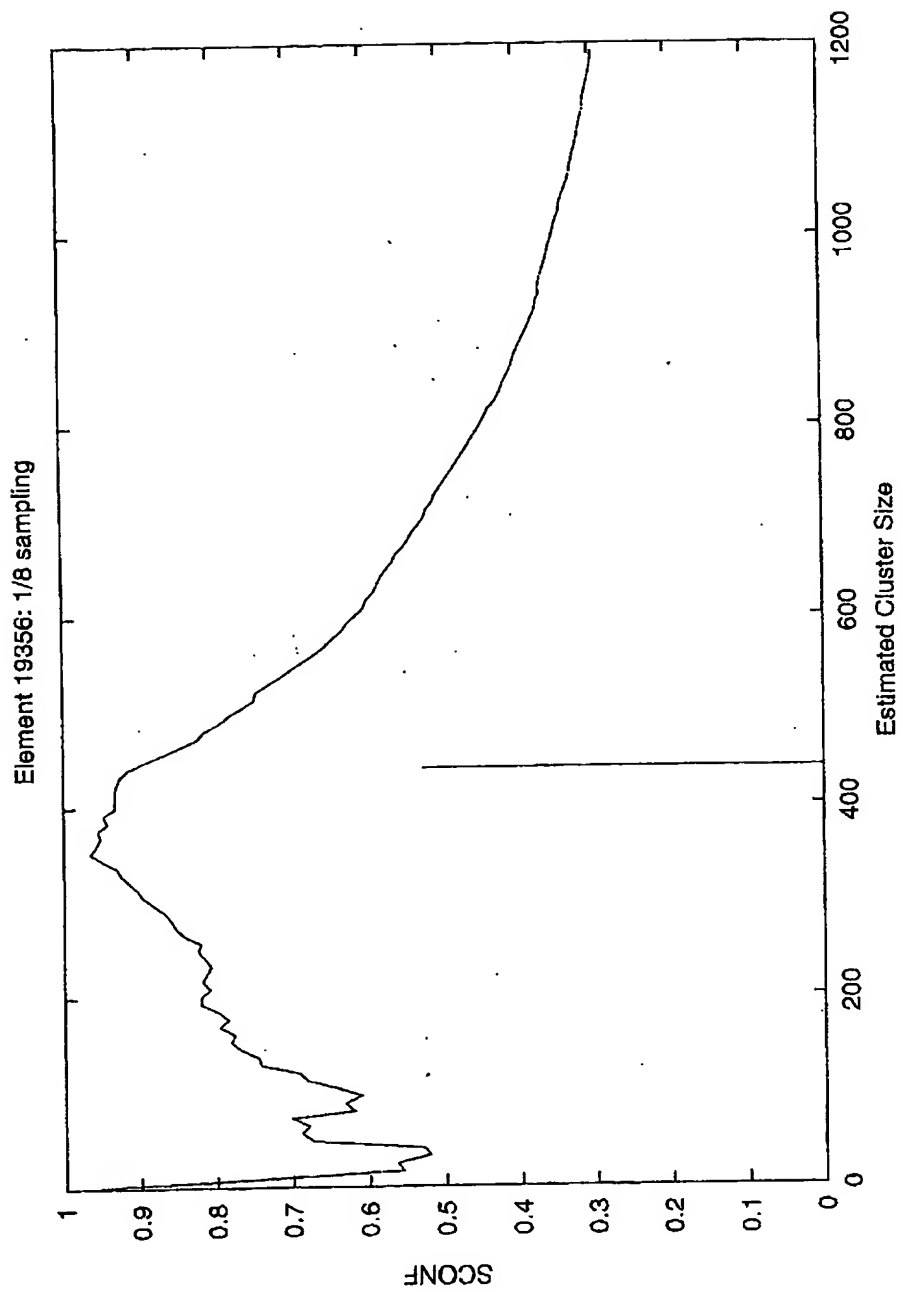


FIG 21